## Second unit: Algorithms

Some topics in Decision Maths; mostly "easy" points on exam.
Largely about optimization problems:

- Kruskal and Prim's algorithms for cheapest spanning trees
- Dijkstra's algorithm: shortest path between two points
- Travelling Salesperson problem

In the real world: computers run these algorithm
From pure math perspective, interesting bits are:

- Proving the algorithm works as advertized
- Analyzing speed of algorithm - can you go faster?

First topic: Prüfer code
How many trees on $n$ vertices are there?

## Two interpretations of counting trees

Count isomorphism classes of trees. "unlabelled trees"

- This is what we do when we count isomers
- No nice answer

Count labelled trees on $n$ vertices

- Vertices are no longer interchangeable
- $n$ ! ways to label an unlabeled tree
- Symmetries mean some produce the same labelled tree

| n | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Unlabelled trees | 1 | 1 | 1 | 2 | 3 | 6 | 11 |
| Labelled trees | 1 | 1 | 3 | 16 | 125 | 1296 | 16807 |

## "Cayley's" formula

Theorem (Borchardt 1860, Cayley 1889)
There are $n^{n-2}$ labelled trees on $n$ vertices.
Original proof used determinants.
Prüfer code: another way to prove Cayley's formula
Bijection between Trees and Codes

- $T_{n}=\{$ labelled trees on $n$ vertices $\}$
- $C_{n}=\left\{\left(a_{1}, a_{2}, \ldots, a_{n-2}\right): a_{i} \in\{1,2, \ldots, n\}\right\}$
- Build a bijection between $T_{n}$ and $C_{n}$
- $\left|C_{n}\right|=n^{n-2}$

Bijective Proofs
In combinatorics, bijective proofs often give more...

## How to write down a labelled trees?

Record the edges:


Each column is an edge

| 1 | 5 | 2 | 6 | 5 | 7 | 8 | 8 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 5 | 2 | 6 | 3 | 4 | 5 | 9 | 5 |

- Records $2 n-2$ numbers between 1 and $n$
- Many ways of recording same tree

Need to record the edges in a canonical way
"Canonical" means: without arbitrary choices

## Prüfer code: iteratively remove lowest leaf

1. Find lowest leaf $\ell$ of $T$
2. Record edge e connecting it to rest of tree
3. Delete $\ell$ and $e$ to get a simpler tree $T^{\prime}$
4. Repeat process with $T^{\prime}$


| Parent | $\underline{5}$ | $\underline{6}$ | $\underline{5}$ | $\underline{2}$ | $\underline{5}$ | $\underline{5}$ | $\underline{8}$ | 9 |
| ---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Leaf | 1 | 3 | 4 | 6 | 2 | 7 | 5 | 8 |

The underlined numbers form the Prüfer code Non-underlined numbers are a permutation of 1-9. Why?

## To see Prüfer code is a bijection, construct inverse

Given a Prüfer code, how to fill in empty boxes?

| Parent | 5 | 6 | 5 | 2 | 5 | 5 | 8 |  |
| ---: | ---: | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| Leaf |  |  |  |  |  |  |  |  |

Numbers in the Prüfer code were parents

- So the numbers in Prüfer code can't be leaves
- We deleted lowest leaf first
- Thus, first leaf is lowest number not in the Prüfer code

To reconstruct tree from code:

- Find lowest number not used yet that's not in remaining code
- Delete the first column
- Iterate; last two numbers left are the last edge


## Cayley's enrichment: keep track of degrees of vertices

| Parent | $\underline{5}$ | $\underline{6}$ | $\underline{5}$ | $\underline{2}$ | $\underline{5}$ | $\underline{5}$ | $\underline{8}$ | 9 |
| ---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Leaf | 1 | 3 | 4 | 6 | 2 | 7 | 5 | 8 |

Degree of vertex $i=$ number of times $i$ appears in table $=$ number of times $i$ appears in Prüfer code +1

Corollary
The number of labeled trees on $n$ vertices where for each $i$, vertex $i$ has degree $d_{i}$ is:

$$
\frac{(n-2)!}{\prod\left(d_{i}-1\right)!}
$$

