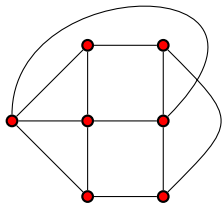
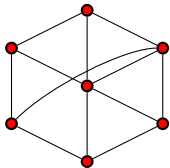


Which graphs are/aren't isomorphic? Prove it.

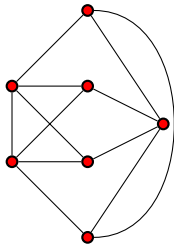
*A*



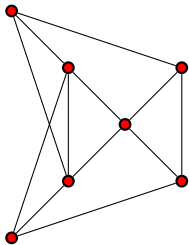
*B*



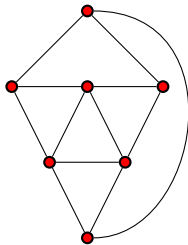
*C*



*D*

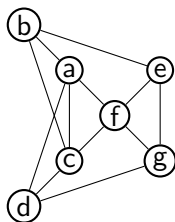
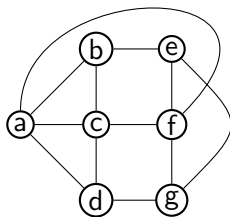


*E*



## One solution to warm-up

- ▶ Graph  $B$  has a vertex of degree 5; others have degree sequence  $[4, 4, 4, 3, 3, 3, 3]$ , so none are isomorphic to  $B$ .
- ▶ In  $A, D, E$ , the three vertices of degree 4 all touch, but not in  $C$ , so none are isomorphic to  $C$ .
- ▶ In  $A, D$ , every vertex is adjacent to a vertex of degree 4, but not in  $E$ , so none are isomorphic to  $E$ .
- ▶ But we see below  $A$  is isomorphic to  $D$ :



# A forest is a bunch of trees

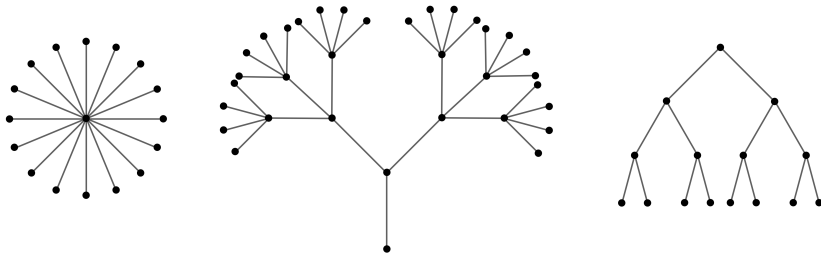


Figure: A forest of three trees

## Definition

- ▶ A *forest* is a graph without cycles
- ▶ A *tree* is a connected graph without cycles

# The Treachery of Definitions (After Magritte)



**Figure:** Ceci n'est pas un arbre (This is not a tree)

## [13/2] ways of looking a tree (After Wallace Stevens)

### Proposition:

Let  $G$  be a graph with  $n$  vertices. The following are equivalent.

1.  $G$  is a tree (i.e.,  $G$  connected but has no cycles)
2. There is a unique path in  $G$  between any two vertices
3.  $G$  is connected and has  $n - 1$  edges
4.  $G$  has no cycles and has  $n - 1$  edges
5.  $G$  is connected, but removing any edge disconnects  $G$
6.  $G$  has no cycles, but adding any edge creates a cycle

### Informally: Trees are Goldilocks graphs

- ▶ Trees have enough edges: they're connected
- ▶ Trees don't have too many edges: they have no cycles

# Make like a tree and get out of here (After Biff Tannen)

## Definition (Tree)

Let  $T$  be a tree. A vertex  $v \in T$  is a *leaf* if it has degree 1.

## Lemma

*Let  $T$  be a tree with  $2 \leq n < \infty$  vertices. Then  $T$  has at least two leaves.*

**Proof 1:** See title of slide.

Pick an edge, and try to “leave” – that is, walk as far as you can.

- ▶ No loops, so you’ll never return to where you are
- ▶ Finitely many vertices, so it can’t go on forever

Eventually you’ll get stuck – that’s a leaf.

# Pruning Trees

## Part of Proposition:

If  $T$  is a tree with  $n$  vertices, then  $T$  has  $n - 1$  edges.

## Proof: Induct on $n$

- ▶ Base case:  $n = 1$
- ▶ Now assume that all trees with  $n - 1$  vertices have  $n - 2$  edges
- ▶ If  $T$  is a tree with  $n$  vertices, it has a leaf  $v$  (by Lemma)
- ▶ Delete  $v$  and the edge next to it to get a new tree  $T'$
- ▶  $T'$  has  $n - 1$  vertices, so  $n - 2$  edges, so  $T$  has  $n - 1$  edges.

## Another use of the handshaking lemma

### Part of Proposition:

If  $G$  is a connected graph with  $n$  vertices and  $n - 1$  edges, then  $G$  is a tree.

### Proof: induct on $n$

- ▶ Base case:  $n = 1$
- ▶ Assume proposition is true for all graphs with  $n - 1$  vertices
- ▶ Since  $G$  is connected, it has no vertices of degree 0
- ▶ Use handshaking to show  $G$  must have a vertex  $v$  of degree 1
- ▶ Delete  $v$  and the edge next to it to get a new graph  $G'$
- ▶  $G'$  is a tree, so  $G$  must have been as well



Chemistry

# Chemical formulas encode degree sequences

| Group→  | 1        | 2        | 3           | 4         | 5         | 6         | 7         | 8         | 9         | 10        | 11        | 12        | 13         | 14        | 15         | 16        | 17         | 18         |         |
|---------|----------|----------|-------------|-----------|-----------|-----------|-----------|-----------|-----------|-----------|-----------|-----------|------------|-----------|------------|-----------|------------|------------|---------|
| ↓Period |          |          |             |           |           |           |           |           |           |           |           |           |            |           |            |           |            |            |         |
| 1       | 1<br>H   |          |             |           |           |           |           |           |           |           |           |           |            |           |            |           |            |            | 2<br>He |
| 2       | 3<br>Li  | 4<br>Be  |             |           |           |           |           |           |           |           |           |           | 5<br>B     | 6<br>C    | 7<br>N     | 8<br>O    | 9<br>F     | 10<br>Ne   |         |
| 3       | 11<br>Na | 12<br>Mg |             |           |           |           |           |           |           |           |           |           | 13<br>Al   | 14<br>Si  | 15<br>P    | 16<br>S   | 17<br>Cl   | 18<br>Ar   |         |
| 4       | 19<br>K  | 20<br>Ca | 21<br>Sc    | 22<br>Ti  | 23<br>V   | 24<br>Cr  | 25<br>Mn  | 26<br>Fe  | 27<br>Co  | 28<br>Ni  | 29<br>Cu  | 30<br>Zn  | 31<br>Ga   | 32<br>Ge  | 33<br>As   | 34<br>Se  | 35<br>Br   | 36<br>Kr   |         |
| 5       | 37<br>Rb | 38<br>Sr | 39<br>Y     | 40<br>Zr  | 41<br>Nb  | 42<br>Mo  | 43<br>Tc  | 44<br>Ru  | 45<br>Rh  | 46<br>Pd  | 47<br>Ag  | 48<br>Cd  | 49<br>In   | 50<br>Sn  | 51<br>Sb   | 52<br>Te  | 53<br>I    | 54<br>Xe   |         |
| 6       | 55<br>Cs | 56<br>Ba | * 71<br>Lu  | 72<br>Hf  | 73<br>Ta  | 74<br>W   | 75<br>Re  | 76<br>Os  | 77<br>Ir  | 78<br>Pt  | 79<br>Au  | 80<br>Hg  | 81<br>Tl   | 82<br>Pb  | 83<br>Bi   | 84<br>Po  | 85<br>At   | 86<br>Rn   |         |
| 7       | 87<br>Fr | 88<br>Ra | * 103<br>Lr | 104<br>Rf | 105<br>Db | 106<br>Sg | 107<br>Bh | 108<br>Hs | 109<br>Mt | 110<br>Ds | 111<br>Rg | 112<br>Cn | 113<br>Uut | 114<br>Fl | 115<br>Uup | 116<br>Lv | 117<br>Uus | 118<br>Uuo |         |
|         |          |          | * 57<br>La  | 58<br>Ce  | 59<br>Pr  | 60<br>Nd  | 61<br>Pm  | 62<br>Sm  | 63<br>Eu  | 64<br>Gd  | 65<br>Tb  | 66<br>Dy  | 67<br>Ho   | 68<br>Er  | 69<br>Tm   | 70<br>Yb  |            |            |         |
|         |          |          | * 89<br>Ac  | 90<br>Th  | 91<br>Pa  | 92<br>U   | 93<br>Np  | 94<br>Pu  | 95<br>Am  | 96<br>Cm  | 97<br>Bk  | 98<br>Cf  | 99<br>Es   | 100<br>Fm | 101<br>Md  | 102<br>No |            |            |         |

| Atom   | C | N | O | H |
|--------|---|---|---|---|
| Degree | 4 | 3 | 2 | 1 |

## Shortcuts around Carbon and Hydrogen

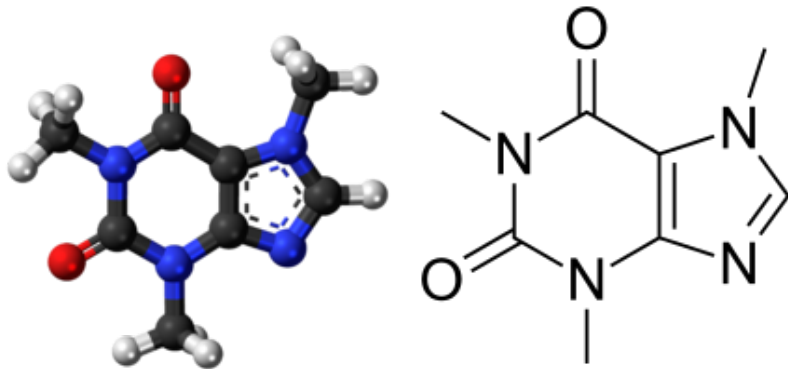


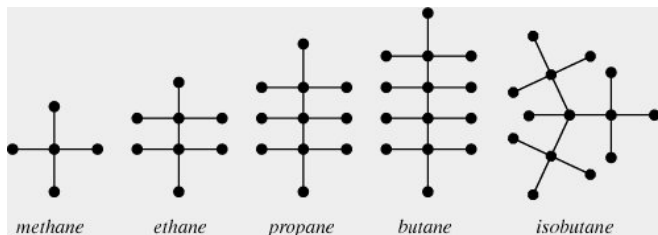
Figure: Two pictures of Caffeine

- ▶ Unlabelled vertices are Carbon
- ▶ Hydrogen not drawn; inferred to make degrees correct

Isomers are graphs with the same degree sequence

### Definition

An *Alkane* is a molecule with formula  $C_nH_{2n+2}$



### Definition (Isomer)

Two different molecules are *isomers* if they have the same chemical formula.

Lemma: Any alkane is a tree.

Proof: Handshaking.

Question: How many isomers does  $C_5H_{12}$  have?